DendrETH Specification

DendrETH is a superset of ALC. Thus, we do not define again the LCU algorithm, but merely the proposed improvements. We define two data structures to use in our protocol. First, the slashing evidence, or evidence. Secondly, we define the slashing action \mathcal{S} . A slashing starts with a beacon state \mathcal{B} [68] and a slashing action. We represent the action of slashing a set of validators \mathcal{N} by $\mathcal{S} \xrightarrow{\text{evidence},\mathcal{N}} \{1\}$

We define evidence as having several attributes:

- 1) BlockHeader*: an attested block header, which is the contender that originates the slashing.
- 2) C_{i+1} : the next sync committee.
- π_{Ci+1}: the next_sync_committee_branch is the Merkle path that authenticates the next sync committee
- 4) BlockHeader: the finalized block header.
- 5) $\pi_{\rm BlockHeader^*}$: a Merkle branch validating the finalized block header.
- 6) aggregate: the aggregated signature $\sigma_{N_{[1.512]}}$ of the nodes participating in the sync committee, $sync_committee_signature$ and the bitmap of the participants bmap.
- 7) *slot*: the signature slot.
- 8) $K_P^{N_1}, \dots, K_P^{N_{512}}$: the sync committee public keys.
- 9) *root*: represents the root hash of a block that the evidence claims to be finalized.
- 10) \(\pi_{root} \): Merkle proof showing the inclusion of \(root \) in the state tree.

It is worth noting that there are two evidences provided in a slashing. The purpose of having two distinct shreds of evidence is to provide proof that a particular validator (or set of validators) committed an equivocation or another malicious act. Equivocation essentially means producing multiple conflicting pieces of information for the same context.

The slashing action S has:

- 1) \mathcal{N} : a list of validators to be slashed.
- 2) evidence.
- BlockHeader_{|HeaderChain|}.root: the recent finalized block header root.
- 4) *slot*: recent finalized slot.

The complete slashing algorithm has two parts: IdentifySlashing and EnforceSlashing. The latter should only be called if the former returns 1 (otherwise, it means that the slashing evidence is not valid).

```
Algorithm 2: Slashing algorithm in DendrETH -
 IdenfitySlashing
  Input: \mathcal{B}, \mathcal{S}, evidence[2]
  Output: \{\bot, 1\}
1 \text{ slash} \leftarrow 0 \triangleright \text{ slash flag}
2 assert S.slot > evidence.slot
3 ▷ asserts both evidences are sequential
4 if evidence[0].slot =
    evidence[1].slot \land evidence[0].Block{\tt Header}^* =
    evidence[1].BlockHeader* then
      slash \leftarrow 1
6 end if
7 if evidence[0].C_{i+1} \neq \text{evidence}[1].C_{i+1} then
s | slash \leftarrow 1
9 end if
10 linear[0] \leftarrow evidence[0].root ==
    evidence[0].BlockHeader.root
11 if \neg(final(evidence[0]) \lor final(evidence[1])) then
      assert BlockHeader<sub>|HeaderChain|</sub>.root = \emptyset
13 end if
14 ▷ Checks to prevent slashing validators
     who signed an alternate history
     non-maliciously
15 canonical_is_0 ←
    evidence[0].BlockHeader.slot \ge
    evidence[1].BlockHeader.slot \land slot ==
    evidence[0].BlockHeader.slot \land
    BlockHeader_{|HeaderChain|}.root ==
    evidence[0].root \land linear[0]
16 ▷ Might check {\cal B} and {\cal S}
17 if canonical is 0 \wedge slash then
18 return 1
19 end if
```

Description of Algorithms 2 and 3 is in the text

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Algorithm 3: Slashing algorithm in DendrETH -
EnforceSlashing
  Input: \mathcal{B}, \mathcal{S}, evidence[2]
  Output: \{\bot, 1\}
1 ToSlash ← \emptyset
2 for key in S.evidence[0].K_P^{N_1}, \ldots, K_P^{N_{512}} do
      ToSlash =
       ToSlash \cup GetPublicKey(key, aggregate.bmap)
4 end for
5 for val in S.N do
     	riangleright asserts at least one validator val to
        be slashed \in \mathcal{B}.validators
7 end for
8 assert ValidateSlashingEvidence(evidence[0],
   S.BlockHeader_{|HeaderChain|}.root, slot,
   \mathcal{B}.GenesisRoot)
9 S \xrightarrow{\text{evidence}, \mathcal{N}} \{1\}
```

Fig. 5: DendrETH specification (slashing algorithm)